# Efficient and Scalable Subgraph Enumeration

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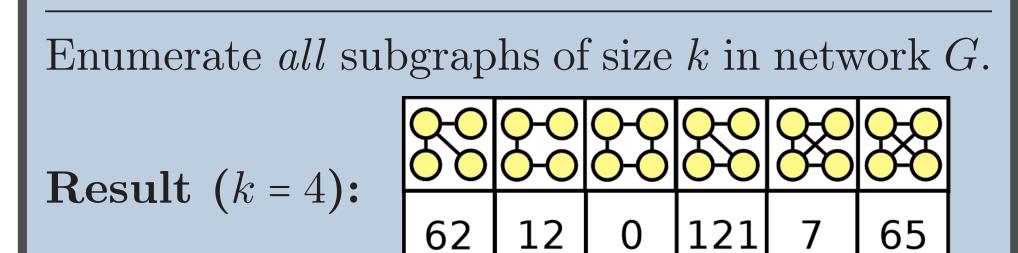


# **Subgraph Enumeration**

#### Summarize/compare networks:

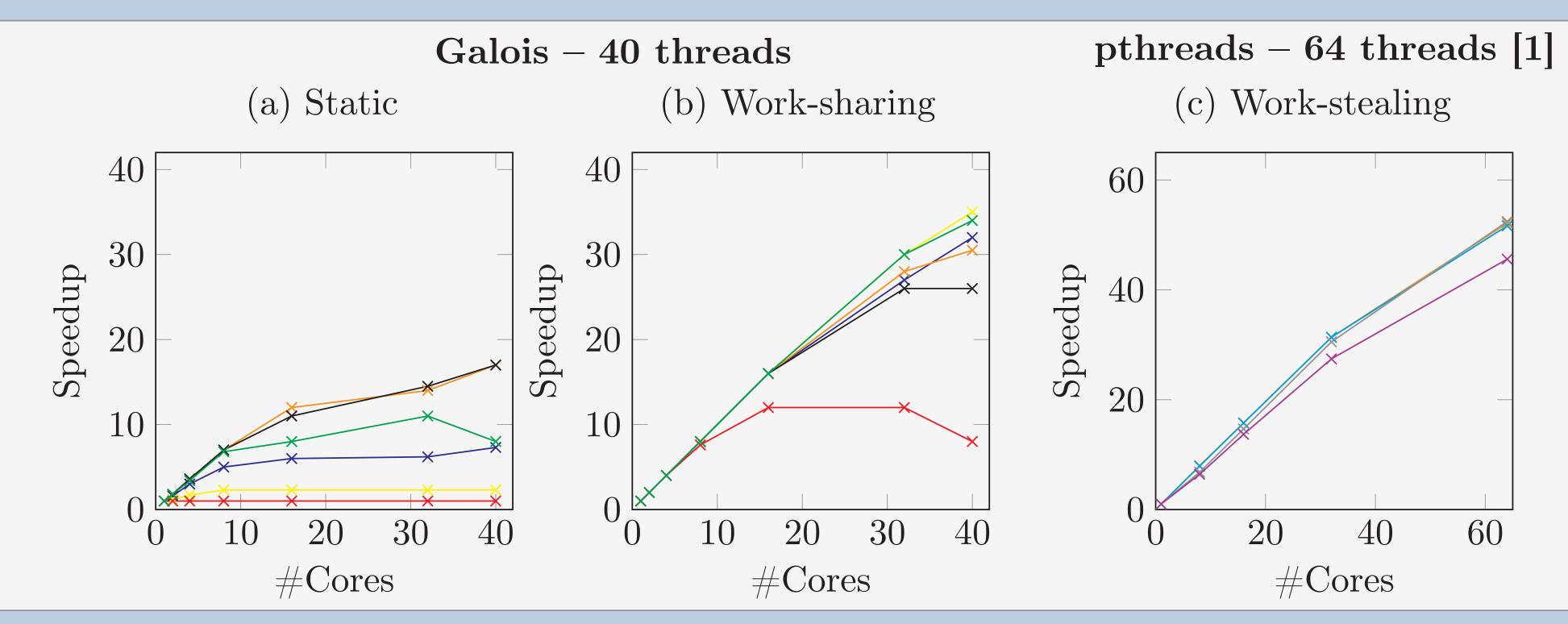
- Centrality measures (closeness, betweenness, PageRank, etc.).
- Degree Distributions/Power law exponent.

• Subgraphs.



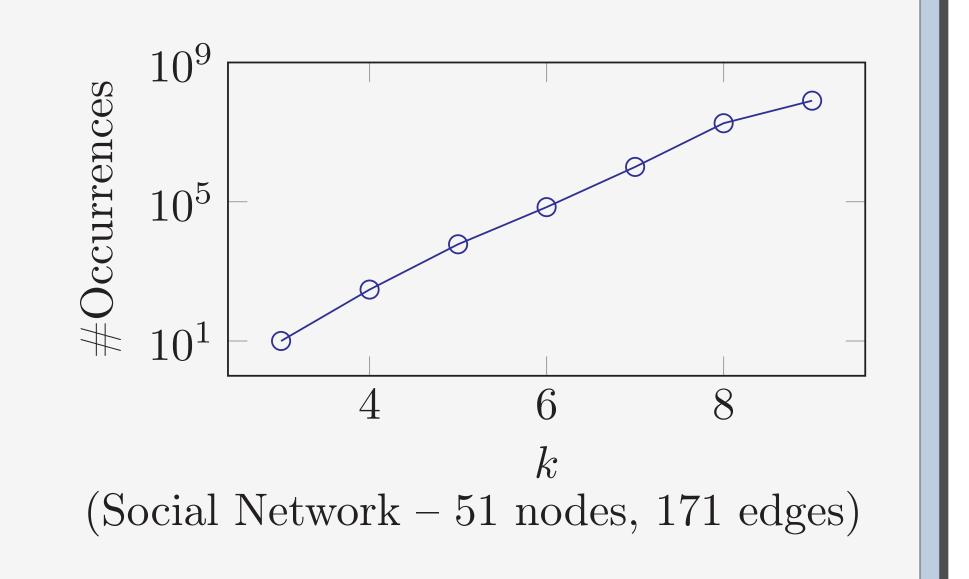
# Parallel Subgraph Enumeration in Shared Memory

Networks: social (jazz, blogs, e-mails), co-authorships (netscience, geometry), biological (neuronal, PPI), linguistic (dictionary) and geometric (routes).



## **Computational Complexity**

• Millions/billions of occurrences which grows exponentially with k.



• Subgraph-types also grow exponentially. (Problematic if algorithm is subgraph-centric)



(a) Achieves some speedup but not for every network. Work is not balanced.

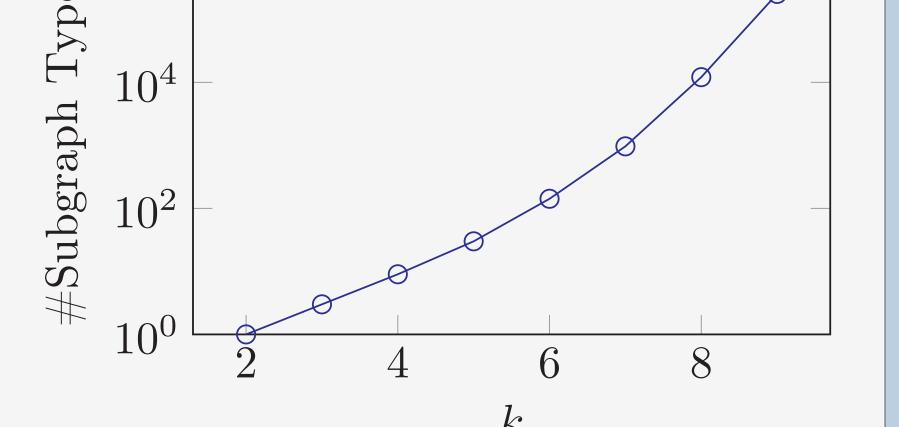
- (b) Flexible but leads to overhead updating the work-sharing queue.
- (c) **Near-linear speedup** for all tested networks.

**Example:** Size-6 census on **blogs** takes  $\approx 2$  days sequentially but only  $\approx 1$  hour using (c).

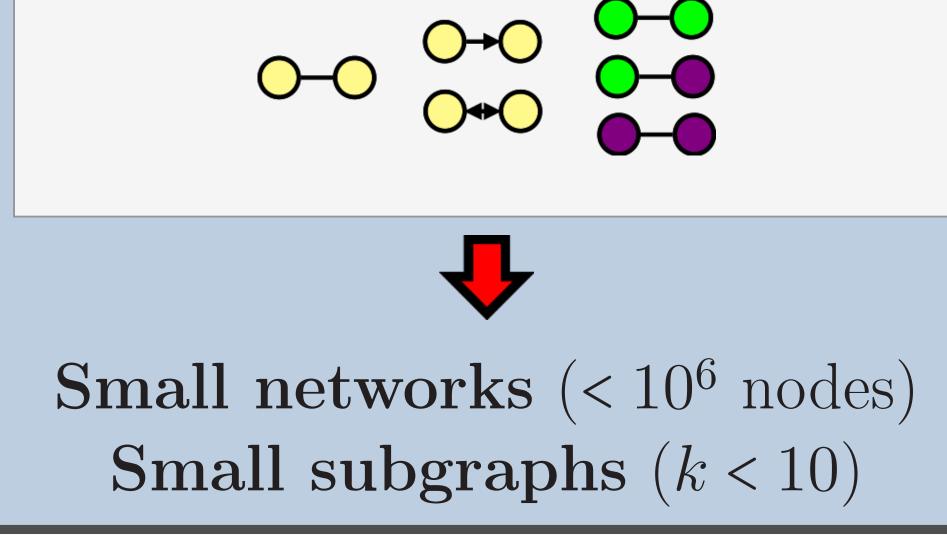
# Distributed Memory, GPU, MapReduce

- **Distributed Memory:** near-linear speedup up to 128 processors.
- **GPU:** hard to distribute work dynamically and efficiently perform graph traversal.
- MapReduce: similar problems to the GPU (ongoing work).

Application: Directed Graphlets [3]



• Edge Direction, Node color, ..., increase the complexity.

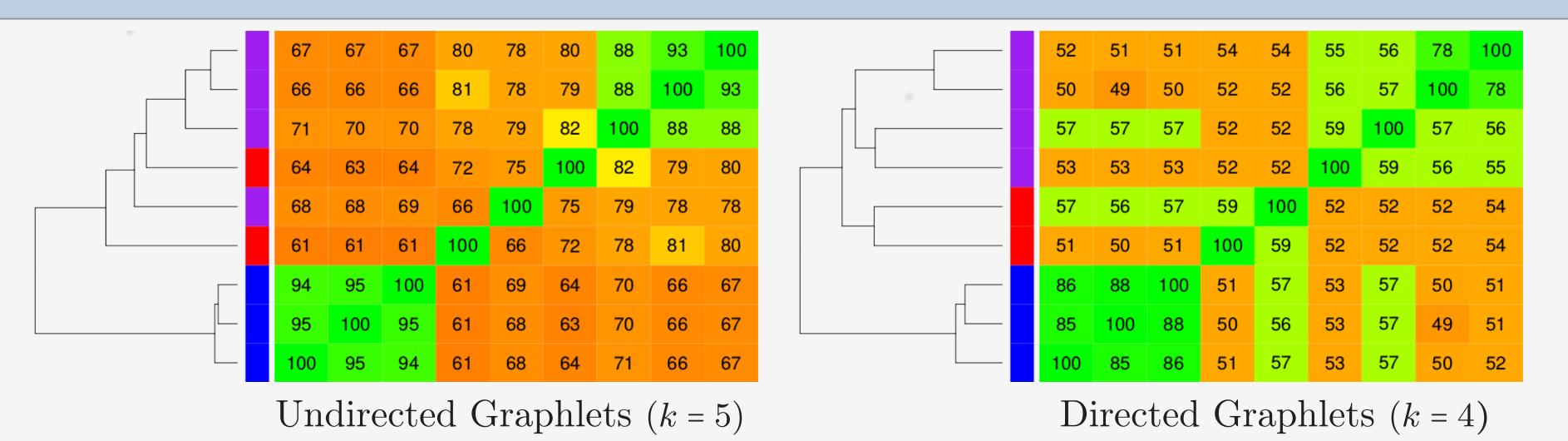


- Node/Network comparison.
- Network Alignment.
- Subgraph-based metrics.



k	2	3	4	5	6	7	8	9
$ u\mathcal{G}_k $	$10^{0}$	$10^{0}$	$10^{1}$	$10^{1}$	$10^{2}$	$10^{3}$	$10^{4}$	$10^{5}$
$egin{array}{c} k \  u {\cal G}_k  \  d {\cal G}_k  \end{array}$	$10^{0}$	$10^{1}$	$10^{2}$	$10^{3}$	$10^{6}$	$10^{8}$	$10^{12}$	$10^{17}$

Directed Biological Networks: cell signaling, metabolic and transcription regulation.



Directed graphlets retrieve relevant topological information

	$\mathcal{G}$	GraphCrunch	Orca	Kavosh ESU		
Speedup:	$\frac{u\mathcal{G}_5}{u\mathcal{G}_6}$	$7.15 \pm 2.56$ n/a	$\begin{array}{r} 2.04 \pm 1.27 \\ n/a \end{array}$	$95.00 \pm 30.97$ $85.89 \pm 24.07$	$\frac{80.11 \pm 27.85}{70.31 \pm 19.88}$	
	$\frac{d\mathcal{G}_4}{d\mathcal{G}_5}$	n/a n/a	$\frac{n/a}{n/a}$	$20.61 \pm 3.80$ $35.00 \pm 9.77$	$18.73 \pm 3.86$ $31.75 \pm 8.30$	

## Our Approach

- 1. Fast sequential algorithm (G-Tries).
- 2. Scalable parallel strategy (Dynamic workload balancing) [1,2].

Our tool is more general and more efficient

### References

[1] D. Aparicio, P. Ribeiro, F. Silva. *Parallel subgraph* counting for multicore architectures in "Parallel and Distributed Processing with Applications", IEEE, 2014.

[2] D. Aparicio, P. Paredes, P. Ribeiro. A Scalable Parallel Approach for Subgraph Census Computation in "Euro-Par 2014: Parallel Processing Workshops", Springer, 2014.

[3] D. Aparicio, P. Ribeiro, F. Silva. "Network comparison using directed graphlets" in arXiv:1511.01964, 2015.

#### Future Work

- **Temporal Graphlets:** compare/summarize temporal networks.
- Subgraph Isomorphism on Streaming Graphs: discover & report blacklisted patterns (colab w/UT-Austin) UTAustin | Portugal International Collaboratory For Emercing Technologies, Collab
- Large Scale Subgraph Enumeration: enumerate bigger subgraphs (k >> 10) on very large networks  $(> 10^9 \text{ nodes}) \rightarrow$  large scale parallel approach.